

Commuting Datatypes

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Generic Programming Summer School
Oxford, August 2002

Introductory Examples

Zip (of lists)

$$([a_1, a_2, \dots, a_n], [b_1, b_2, \dots, b_n]) \mapsto [(a_1, b_1), (a_2, b_2), \dots, (a_n, b_n)]$$

$$\text{Pair} \cdot \text{List} \quad \mapsto \quad \text{List} \cdot \text{Pair}$$

Matrix Transposition

$$\text{List} \cdot \text{List} \quad \mapsto \quad \text{List} \cdot \text{List}$$

Broadcast

$$(a, [b_1, b_2, \dots, b_n]) \mapsto [(a, b_1), (a, b_2), \dots, (a, b_n)]$$

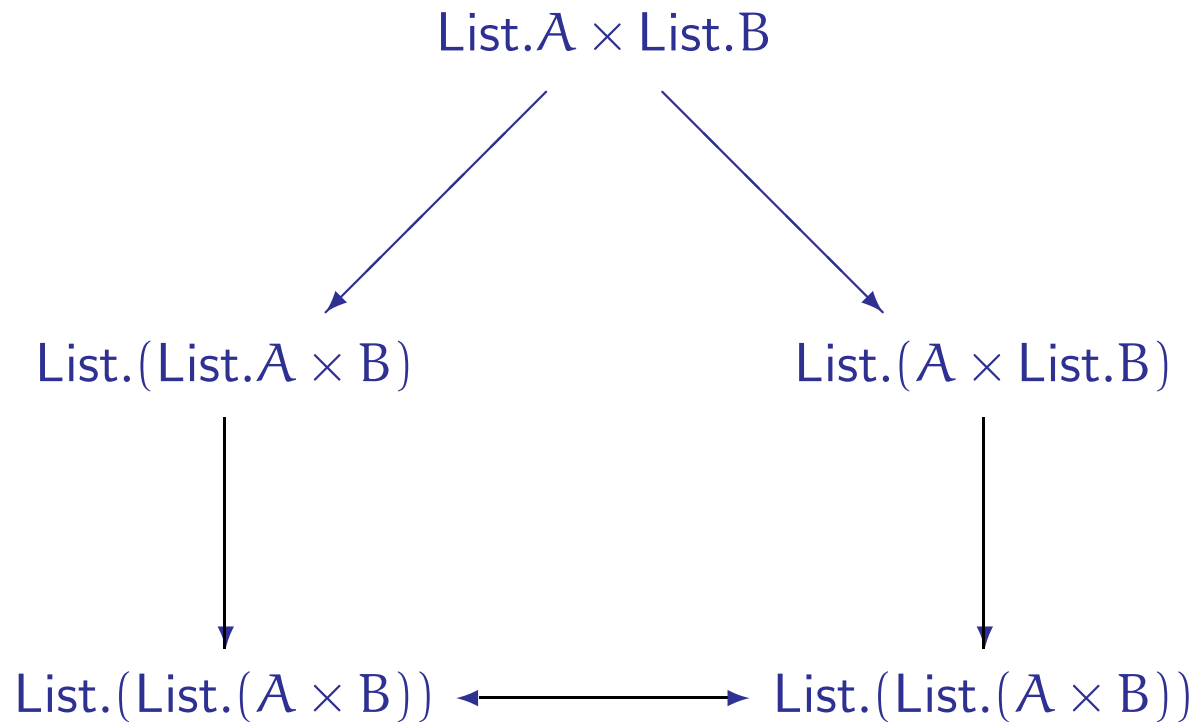
$$A \times \cdot \text{List} \quad \mapsto \quad \text{List} \cdot A \times$$

Primitive

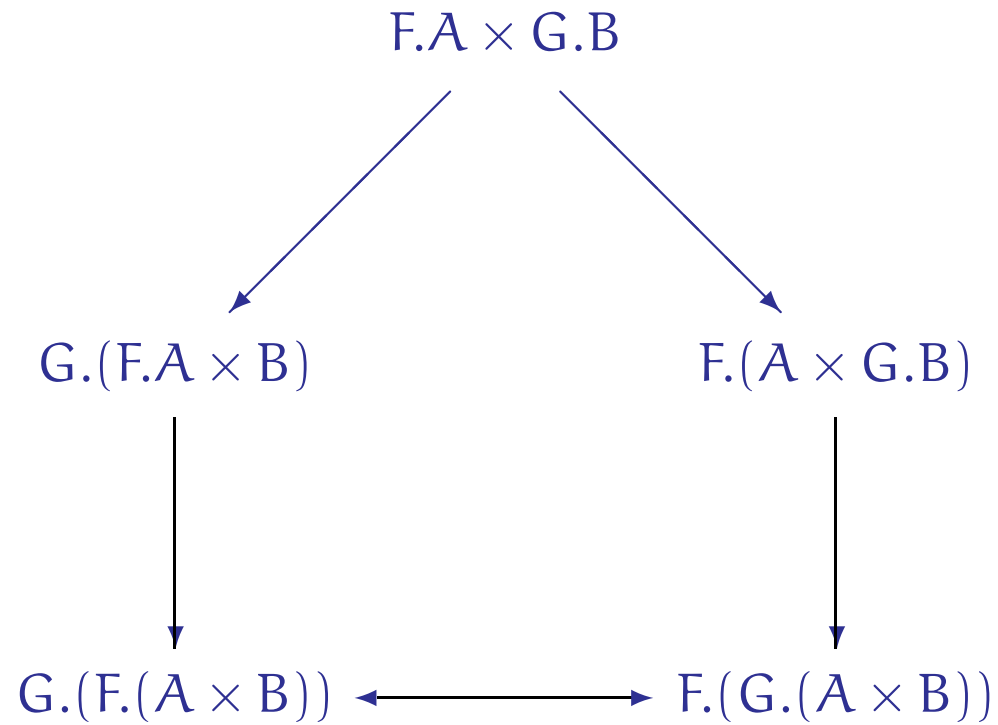
$$(A+B) \times (C+D) \mapsto (A \times C) + (B \times D)$$

$$\times \cdot + \quad \mapsto \quad + \cdot \times$$

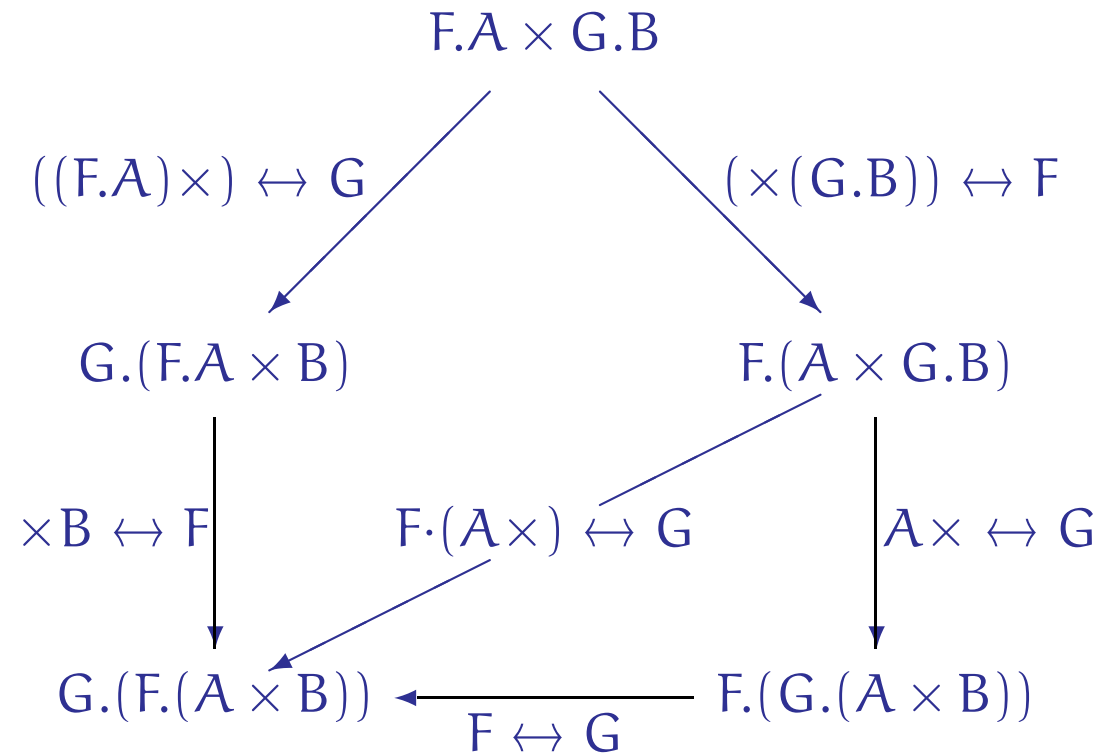
Structure Multiplication ...



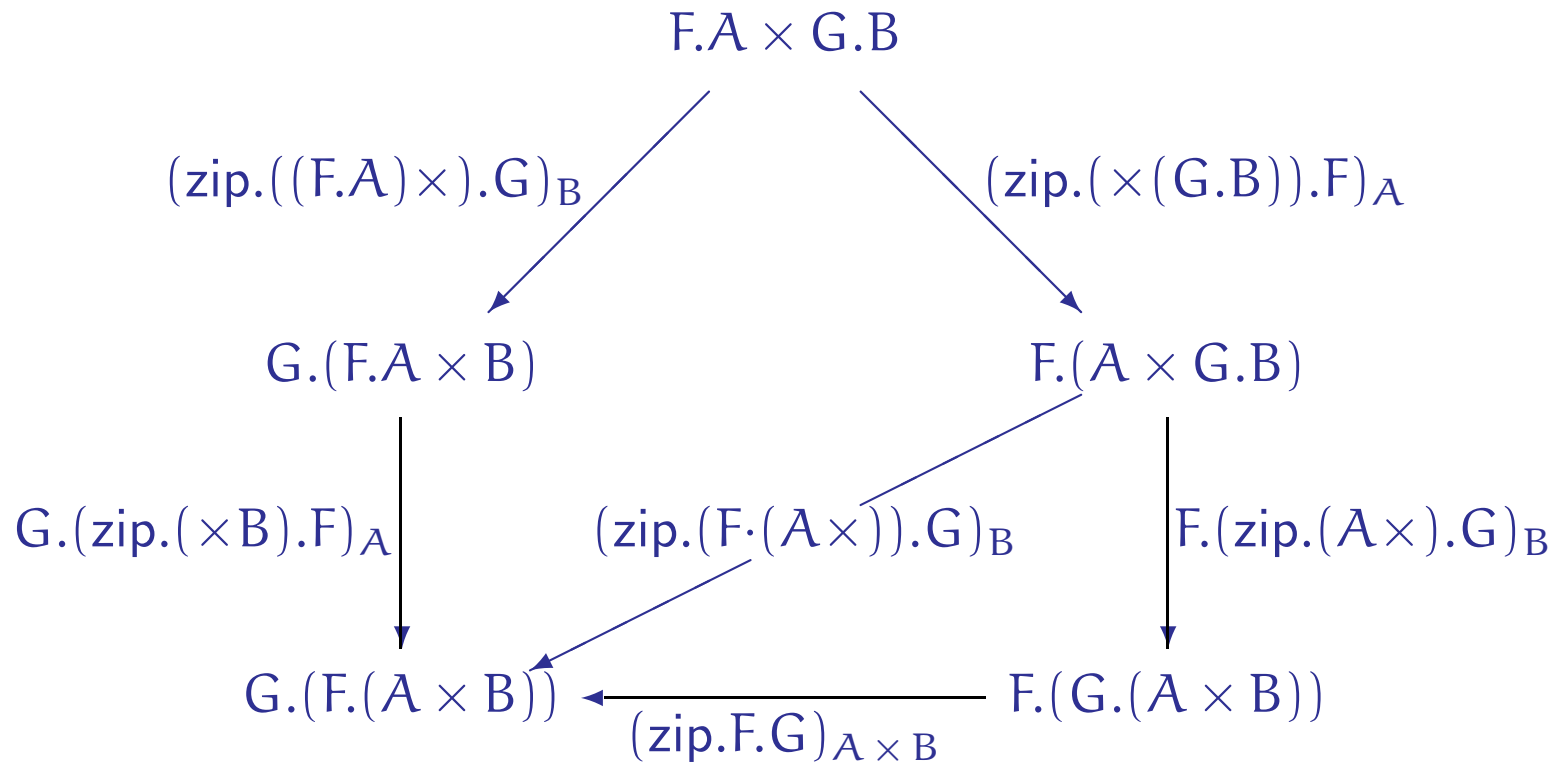
... Generalised ...



... Illustrates Generic Requirements



Multi-Coloured Zips



Broadcasts ...

A broadcast copies a given value across all storage locations of a datatype.

Formally, a family of functions bcst , where

$$\text{bcst}_{A,B} : F.(A \times B) \leftarrow F.A \times B$$

is said to be a *broadcast* for datatype F iff it is parametrically polymorphic in the parameters A and B and $\text{bcst}_{A,B}$ behaves coherently with respect to product in the following sense:

... Respect the Unit of Product ...

The following diagram

$$\begin{array}{ccc}
 F.(A \times \mathbb{1}) & \xleftarrow{\text{bcst}_{A, \mathbb{1}}} & F.A \times \mathbb{1} \\
 \searrow (F \cdot \text{rid})_A & & \swarrow (\text{rid} \cdot F)_A \\
 & F.A &
 \end{array}$$

(where $\text{rid}_A : A \leftarrow A \times \mathbb{1}$ is the obvious natural isomorphism) commutes.

... and Associativity of Product

The following diagram

$$\begin{array}{ccc}
 F.A \times (B \times C) & \xleftarrow{\text{ass}_{F.A, B, C}} & (F.A \times B) \times C \\
 & & \downarrow \text{bcst}_{A, B} \times \text{id}_C \\
 & & F.(A \times B) \times C \\
 & & \downarrow \text{bcst}_{A \times B, C} \\
 & & F.((A \times B) \times C) \\
 & \xleftarrow{F \cdot \text{ass}_{A, B, C}} & \\
 F.(A \times (B \times C)) & &
 \end{array}$$

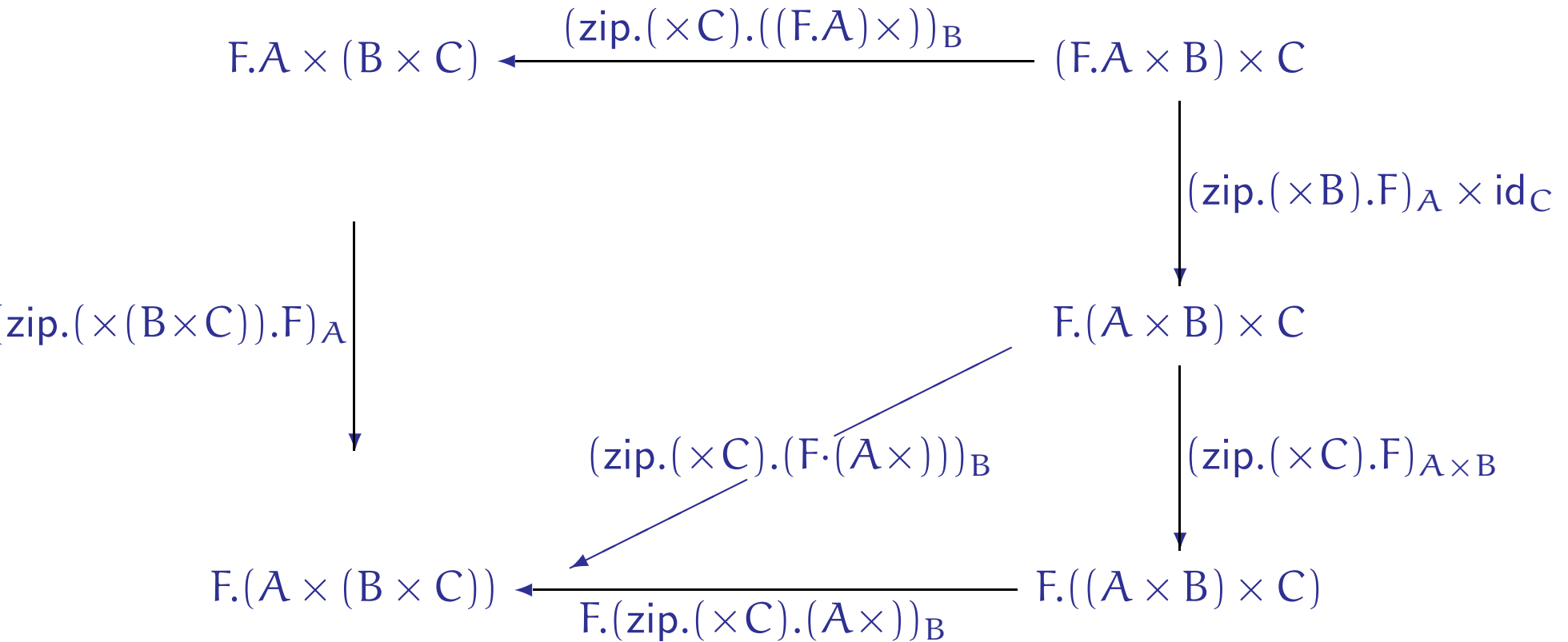
$\text{bcst}_{A, B \times C}$ (vertical arrow from $F.A \times (B \times C)$ to $F.(A \times (B \times C))$)

(where $\text{ass}_{A, B, C} : A \times (B \times C) \leftarrow (A \times B) \times C$ is the obvious natural isomorphism) commutes as well.

Unit of Product is a “zip”

$$\begin{array}{ccc}
 F.(A \times \mathbb{1}) & \xleftarrow{(\text{zip} . (\times \mathbb{1}) . F) \cdot K_A} & F.A \times \mathbb{1} \\
 \searrow & & \swarrow \\
 F.(\text{zip} . (\times \mathbb{1}) . (K_A)) & & \text{zip} . (\times \mathbb{1}) . (K_{F.A}) \\
 & \searrow & \swarrow \\
 & F.A &
 \end{array}$$

Associativity of Product is a “Zip”



Conclusion

- Commuting Datatypes (“Zips”) are everywhere!
- Generic specification and proof is (potentially) very effective.
- A relational framework is necessary.
- Challenge: give generic specification of “commuting datatypes” from which “zips” can be constructed calculationaly.