

Computer Systems Architecture

<http://cs.nott.ac.uk/~txa/g51csa/>

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Lecture 05: Comparisons, Loops and Bitwise Operations



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Using C

- Java syntax is based on C.
- C is more low-level than Java:
 - Pointers.
 - goto.
- C++ is an extension of C.
- wikibook: *Programming in C* available at http://en.wikibooks.org/wiki/C_Programming



Hello World in C

```
#include <stdio.h>

int main() {
    printf("Hello, world!\n");
}
```

- Store in `hello.c`.
- Compile with:
`gcc hello.c -o hello`
- Under UNIX run with:
`hello`



Inequalities

- Previously we learnt beq and bne (branch on = and \neq)
 - Can implement `if(a == b) ...` and `if(a != b) ...`
- But we want other arithmetic comparison operators:

Operator	Name	Abbreviation
=	<i>equals</i>	eq
\neq	<i>not equals</i>	ne
<	<i>less than</i>	lt
\leq	<i>less than or equals</i>	le
>	<i>greater than</i>	gt
\geq	<i>greater than or equals</i>	ge



Comparison Instructions

`slt dst, src0, src1` – Set on Less Than

- Set `dst` to 1 if `src0` is less than `src1`, otherwise 0
 - `if(src0 < src1)`
 - `dst := 1;`
 - `else`
 - `dst := 0;`
 - `dst := src0 < src1 ? 1 : 0;`
-
- There is also `slti` – Set on Less Than Immediate
 - Other pseudoinstructions: `seq`, `sne`, `sle`, `sgt`, `sge`, ...



Decisions on Inequalities

- How do we implement `if(a < b) { ... }`,
given `if(c != d) { ... }` and `c = a < b ? 1 : 0` ?
- We can use two comparisons:
`c = a < b ? 1 : 0;`
`if(c != 0) { ... }`
- Suppose `a` and `b` are `$s0` and `$s1` respectively:
`slt $t0, $s0, $s1`
`beq $t0, $zero, a_ge_b`
`# then-block`
`a_ge_b: # rest of program`



Example: Maximum of Two Numbers

- Given two numbers x and y , calculate which is the larger and store it in m
- In Java/C: $m = x$; $\text{if}(m < y) m = y$;
- In MIPS assembly, with $\$s0$, $\$s1$ and $\$a0$ for x , y and m :

```
move $a0, $s0
slt $t0, $a0, $s1
beq $t0, $zero, a0_ge_s1
move $a0, $s1
```

```
a0_ge_s1:
```

```
# Largest number now in $a0
```



A Note on Pseudoinstructions

- The MIPS *processor* only has `slt`, `beq` and `bne` ...
- But the *assembler* also accepts `sge`, `blt`, `ble` and so on
 - These are *pseudoinstructions* (like `li`, `move`, ...)
- Pseudoinstructions makes assembly programming easier
 - Write what we mean, not what the processor can do
 - Let the assembler insert the necessary instructions
 - Assembler uses `$at` to implement pseudoinstructions
- In the previous example, we can replace:

```
slt $t0, $a0, $s1  
beq $t0, $zero, a0_ge_s1
```

with:

```
bge $a0, $s1, a0_ge_s1
```



While Loops

- Loops are important building blocks in larger programs
- `while` repeats code block as long as condition holds
- What if $i \geq 8$ before the loop begins?

Java/C

```
while(i < 8) {  
    j = j + 3;  
    i = i + 1;  
}
```

C (using goto)

```
    goto while_cond;  
while_loop:  
    j = j + 3;  
    i = i + 1;  
while_cond:  
    if(i < 8) goto while_loop;
```

Implementing While Loops

- Rewritten in C using labels and gotos
 - Made our high-level description more concrete
 - Easier to read than assembly instructions
 - Each line has a simple and direct MIPS implementation

- Assume

i	j
\$s0	\$s1

```

        j while_cond
while_loop:
        addi $s1, $s1, 3
        addi $s0, $s0, 1
while_cond:
        blt $s0, 8, while_loop
  
```



For Loops

- For loops consist of initialiser, condition and counter parts
- A for loop is just a syntactic shortcut for a while loop
 - ... but we already know how to implement while loops!

For Loop

```
for(i = 0; i < 8; i = i + 1)
    j = j + 3;
```

While Equivalent

```
i = 0;
while(i < 8) {
    j = j + 3;
    i = i + 1;
}
```



Example: How Long is a String?

- Arrive at assembly via a series of translations
- *p means “look up the contents of memory location p”

C (For Loop)

```
length = 0;
```

```
for(p = string; *p != 0; p++)  
    length++;
```

Example: How Long is a String?

- Arrive at assembly via a series of translations
- *p means “look up the contents of memory location p”

C (While Loop)

```
length = 0;
p = string;

while(*p != 0) {
    length++;
    p++;
}
```

Example: How Long is a String?

- Arrive at assembly via a series of translations
- *p means “look up the contents of memory location p”

C (using goto)

```
length = 0;
p = string;
goto strlen_cond
strlen_loop:
    length++;
    p++;
strlen_cond:
    c = *p;
    if(c != 0) goto strlen_loop;
```

Example: How Long is a String?

- Arrive at assembly via a series of translations
- `length = $v0, p = $a0, c = $t0`

MIPS Assembly

```
li $v0, 0
la $a0, string
j strlen_cond
strlen_loop:
    addi $v0, $v0, 1
    addi $a0, $a0, 1
strlen_cond:
    lbu $t0, ($a0)
    bne $t0, $zero, strlen_loop
```

Shift to the Left and Shift to the Right

- Shifts move a word's bit pattern to the left or right
- Each shift left ($x = x \ll 1$ in Java syntax)
 - Drops the most significant (leftmost) bit
 - Appends a 0 bit to the least significant end (right)
 - Equivalent to multiplying by 2, ignoring overflow
 - e.g. $0000\ 0101_2 \ll 3 = 0010\ 1000_2$
- Each shift right ($x = x \gg 1$ in Java syntax)
 - Drops the least significant (rightmost) bit
 - Prepends a 0 bit to the most significant end (left)
 - Equivalent to dividing by 2, ignoring remainder
 - e.g. $1001\ 0011_2 \gg 3 = 0001\ 0010_2$



Shift Instructions

`sll dst, src, shamt` – shift left logical

- $dst := src \ll shamt$

`srl dst, src, shamt` – shift right logical

- $dst := src \gg shamt$

Example

Before

```
$s0 = 7C08 02A616  
= 0111 1100 0000 1000 0000 0010 1010 01102
```

```
srl $s0, $s0, 8  
sll $s0, $s0, 12
```

After

```
$s0 = C080 200016  
= 1100 0000 1000 0000 0010 0000 0000 00002
```

Bitwise Logical Operations

- Bitwise – no interaction between different bits of a word
- AND (&) can be used for testing certain bits of a word
- OR (|) can be used for setting certain bits of a word
- XOR (^) can be used for inverting certain bits of a word
- NOT (~) inverts all the bits in a word

a	1100 ₂
b	1010 ₂
a & b	1000 ₂

a	1100 ₂
b	1010 ₂
a b	1110 ₂

a	1100 ₂
b	1010 ₂
a ^ b	0110 ₂



Bitwise Logical Instructions

- Immediate variants omitted: `andi`, `ori` and `xori`

`and dst, src0, src1` – Bitwise AND

- $dst := src_0 \& src_1$

`or dst, src0, src1` – Bitwise OR

- $dst := src_0 | src_1$

`xor dst, src0, src1` – Bitwise XOR

- $dst := src_0 \hat{ } src_1$

`nor dst, src0, src1` – Bitwise NOR

- $dst := \sim(src_0 | src_1)$
- To get bitwise NOT: `nor dst, src, $zero`