

4th Coursework

23/02/2004

Deadline: 27/02/2004 - 15:30 (B33)

Apply the pumping lemma for regular languages to show that the following languages are not regular:

1. We use $\Sigma_1 = \{a, b, c\}$ and

$$L_1 = \{a^n b^m c^{n+m} \mid m, n \in \text{Nat}\}$$

I.e. $aabbcccc \in L_1$ but $aabbcc \notin L_1$

2. We consider the language of repetitions over $\Sigma_2 = \{0, 1\}$ that is

$$L_2 = \{ww \mid w \in \Sigma_2^*\}$$

I.e. $011011 \in L_2$ (using $w = 011$) but $01010 \notin L_2$ (because it cannot be read as a repetition).

What happens if we consider $\Sigma = \{1\}$ instead?