School for Computer Science and Information Technology Machines and their languages (G51MAL) Spring 2004 Dr. Thorsten Altenkirch

## Solutions to 1st Coursework

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1. (a) Transition diagram for A



	i.	$\epsilon \in L(A)$	$\epsilon \notin L(B)$	
	ii.	$aabb \in L(A)$	aabb $\notin L(B)$	
	iii.	aaab $\notin L(A)$	$aaab \in L(B)$	
	iv.	$bbb \in L(A)$	$bbb \notin L(B)$	
(c)				

$$\hat{\delta}_A(0, \text{bab}) = \hat{\delta}_A(\delta_A(0, \text{b}), \text{ab}) \quad \text{definition of } \hat{\delta} \\ = \hat{\delta}_A(2, \text{ab}) \qquad \text{because } \delta_A(0, \text{b}) = 2 \\ = \hat{\delta}_A(\delta_A(2, \text{a}), \text{b}) \quad \text{definition of } \hat{\delta} \\ = \hat{\delta}_A(3, \text{b}) \qquad \text{because } \delta_A(2, \text{a}) = 3 \\ = \hat{\delta}_A(\delta_A(3, \text{b}), \epsilon) \quad \text{definition of } \hat{\delta} \\ = \hat{\delta}_A(1, \epsilon) \qquad \text{because } \delta_A(3, \text{b}) = 1 \\ = 1 \qquad \text{definition of } \hat{\delta}$$

$$\begin{split} \hat{\delta}_B(0, \text{bab}) &= \hat{\delta}_B(\delta_B(0, \text{b}), \text{ab}) & \text{definition of } \hat{\delta} \\ &= \hat{\delta}_B(0, \text{ab}) & \text{because } \delta_B(0, \text{b}) = 0 \\ &= \hat{\delta}_B(\delta_B(0, \text{a}), \text{b}) & \text{definition of } \hat{\delta} \\ &= \hat{\delta}_B(1, \text{b}) & \text{because } \delta_B(0, \text{a}) = 1 \\ &= \hat{\delta}_B(\delta_B(1, \text{b}), \epsilon) & \text{definition of } \hat{\delta} \\ &= \hat{\delta}_B(3, \epsilon) & \text{because } \delta_B(1, \text{b}) = 3 \\ &= 3 & \text{definition of } \hat{\delta} \end{split}$$

(d) L(A) contains all the words such that the number of a's and b's have a different remainder when divided by 2. Writing #(x, w) for the number of x's in w we can express this by:

 $L(A) = \{ w \mid \#(a, w) \not\equiv \#(b, w) \mod 2 \}$ 

 ${\cal L}(B)$  contains all words such that the letter before the last one is a.

 $L(B) = \{ wax \mid w \in \{a, b\}^*, x \in \{a, b\} \}$ 

2.

