

# Exercises for Naive Type Theory

## Part 3

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- Given a type  $A$  and  $n : \mathbb{N}$  we define the type  $\text{Vec } A n$  as given by the constructors:

$$\begin{aligned} [] &: \text{Vec } A 0 \\ - :: - &: \Pi_{n:\mathbb{N}} A \rightarrow \text{Vec } A n \rightarrow \text{Vec } A (1 + n) \end{aligned}$$

We define lists over  $A$  the following way:

$$\text{List } A := \Sigma n : \mathbb{N}. \text{Vec } A n$$

Given this derive the following operations on lists:

$$\begin{aligned} \text{nil} &: \text{List } A \\ \text{cons} &: A \rightarrow \text{List } A \rightarrow \text{List } A \\ \text{fold} &: X \rightarrow (A \rightarrow X \rightarrow X) \rightarrow \text{List } A \rightarrow X \end{aligned}$$

using general recursive definitions by pattern matching.

- Given  $n : \mathbb{N}$  we define the type  $\text{Fin } n$  as given by the following constructors

$$\begin{aligned} 0 &: \Pi_{n:\mathbb{N}} \text{Fin } 1 + n \\ \text{suc} &: \Pi_{n:\mathbb{N}} \text{Fin } n \rightarrow \text{Fin } (1 + n) \end{aligned}$$

Define the following functions

$$\begin{aligned} \text{max} &: \Pi_{n:\mathbb{N}} \text{Fin}(1 + n) \\ \text{emb} &: \Pi_{n:\mathbb{N}} \text{Fin } n \rightarrow \text{Fin } (1 + n) \\ \text{inv} &: \Pi_{n:\mathbb{N}} \text{Fin } n \rightarrow \text{Fin } (1 + n) \end{aligned}$$

where  $\text{max}$  computes the maximal element in a non-empty finite type;  $\text{emb}$  embedds  $\text{Fin } n$  into  $\text{Fin } (1 + n)$  without changing the value, e.g.  $\text{emb } 3_5 \equiv 3_6$ . The function  $\text{inv}$  inverts the order of elements in a finite type in particular it maps  $0$  to the maximal element and the maximal element to  $0$ .

3. Using propositions as types, can you prove the translation of the axiom of choice? Let  $A, B$  be type and  $R$  be a relation between  $A$  and  $B$  the axiom of choice is:

$$(\forall x : A. \exists y : B. R x y) \Rightarrow (\exists f : A \rightarrow B. \forall x : A. R x (f x))$$

4. Using propositions as types try to prove the de Morgan laws of predicate logic:

$$\neg(\forall x : A. P x) \Leftrightarrow \exists x : A. \neg(P x)$$

$$\neg(\exists x : A. P x) \Leftrightarrow \forall x : A. \neg(P x)$$