

Exercises for Naive Type Theory

Part 4

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1. Given the definition of *Tree* : **Type** by the following constructor:

$\text{mkTree} : \Pi I : \mathbf{Type}. (I \rightarrow \text{Tree}) \rightarrow \text{Tree}$

Implement the following axioms of set theory as operations on trees:

- (a) Empty set axiom

$$\exists x. \forall y. \neg(y \in x)$$

- (b) Axiom of pairing

$$\forall x, y. \exists z. x \in z \wedge y \in z$$

- (c) Axiom of union

$$\forall x. \exists y. (\forall z. z \in x. \forall w. w \in z \Rightarrow w \in y)$$

- (d) Axiom of infinity

$$\exists x. \emptyset \in x \wedge \forall y. y \in x \Rightarrow y \cup \{y\} \in x$$

- (e) Powerset axiom

$$\forall x. \exists y. \forall z. z \subseteq x \Rightarrow z \in y$$

2. Define the function

$\text{min} : \mathbb{N} \rightarrow \mathbb{N} \rightarrow \mathbb{N}$

which computes the minimum of two natural numbers first using pattern matching style and then only using

$\text{rec}_A : A \rightarrow (\mathbb{N} \rightarrow A \rightarrow A) \rightarrow \mathbb{N} \rightarrow A$

$$\begin{aligned} \text{rec } z \text{ s } 0 &::= z \\ \text{rec } z \text{ s } (\text{suc } m) &::= s \text{ n } (\text{rec } z \text{ s } m) \end{aligned}$$

Prove that

$$\forall x, y : \mathbb{N}. \text{min } x \text{ y} = \text{min } y \text{ x}$$

using propositions as types. First give a proof using pattern matching style, then give a proof only using the eliminator:

$$\text{elim} : \Pi A : \mathbb{N} \rightarrow \mathbf{Type}. A 0 \rightarrow (\Pi n : \mathbb{N}. A n \rightarrow A (\text{suc } n)) \rightarrow \Pi n : \mathbb{N}. A n$$
$$\begin{aligned} \text{elim } A \text{ z } s \text{ 0} &::= z \\ \text{elim } A \text{ z } s (\text{suc } n) &::= s \text{ n } (\text{elim } A \text{ z } s \text{ n}) \end{aligned}$$

3. We define the dependent type

$$\text{In} : A \rightarrow \text{List } A \rightarrow \mathbf{Type}$$

as given by the constructors:

$$\begin{aligned} \text{here} &: \Pi a : A. \Pi l : \text{List } A. \text{In } a (a :: l) \\ \text{later} &: \Pi a, b : A. \Pi l : \text{List } A. \text{In } a \text{ l} \rightarrow \text{In } a (b :: l) \end{aligned}$$

Derive the recursor (non dependent eliminator) and the dependent eliminator for this In .