Naturality for Free

The category interpretation of directed type theory

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reverse is natural

$$\mathrm{List} : \underline{\mathbf{Set}} \to \underline{\mathbf{Set}}$$
$$\mathrm{rev} : \Pi_{A:\mathbf{Set}} \mathrm{List}\, A \to \mathrm{List}\, A$$
$$f : A \to B$$

$$\begin{array}{ccc}
\mathrm{List}\, A & \xrightarrow{\mathrm{rev}_A} & \mathrm{List}\, A \\
{\scriptstyle \mathrm{List}\, f}\downarrow & & \downarrow{\scriptstyle \mathrm{List}\, f} \\
\mathrm{List}\, B & \xrightarrow[\mathrm{rev}_B]{} & \mathrm{List}\, B
\end{array}$$

Proof by ...-induction

$$\mathrm{List}\, f\, [a_0, a_1, \dots, a_{n-1}] = [f\ a_0, f\ a_1, \dots, f\ a_{n-1}]$$

$$\mathrm{rev}_A\, [a_0, a_1, \dots, a_{n-1}] = [a_{n-1}, \dots, a_1, a_0]$$

$$\begin{aligned}
(\mathrm{rev}_B \circ \mathrm{List}\, f)\, [a_0, a_1, \dots, a_{n-1}] &= \mathrm{rev}_B\, (\mathrm{List}\, f\, [a_0, a_1, \dots, a_{n-1}]) \\
&= \mathrm{rev}_B\, [f\ a_0, f\ a_1, \dots, f\ a_{n-1}]) \\
&= [f\ a_{n-1}, \dots, f\ a_1, f\ a_0]) \\
&= \mathrm{List}\, f\, [a_{n-1}, \dots, a_1, a_0] \\
&= \mathrm{List}\, f\, (\mathrm{rev}_A\, [a_0, a_1, \dots, a_{n-1}]) \\
&= (\mathrm{List}\, f \circ \mathrm{rev}_A)\, [a_0, a_1, \dots, a_{n-1}]
\end{aligned}$$

Everything is natural ...

$$F, G : \underline{\mathbf{Set}} \to \underline{\mathbf{Set}}$$

$$\alpha : \Pi_{A:\mathbf{Set}} F\,A \to G\,A$$

$$f : A \to B$$

$$\begin{CD}
F\,A @>{\alpha_A}>> G\,A \\
@V{F\,f}VV @VV{G\,f}V \\
F\,B @>>{\alpha_B}> G\,B
\end{CD}$$

. . . but we can't prove it.

- We know that all families of functions are natural.
- But we cannot prove it.
- It should be a *free theorem*.

The hint (HoTT)

$$F, G : \underline{\mathbf{Set}} \to \underline{\mathbf{Set}}$$
$$\alpha : \Pi_{A:\mathbf{Set}} F\,A \simeq G\,A$$
$$f : A \simeq B$$

$$\begin{array}{ccc} F\,A & \xrightarrow{\alpha_A} & G\,A \\ {\scriptstyle F\,f}\downarrow & & \downarrow{\scriptstyle G\,f} \\ F\,B & \xrightarrow[\alpha_B]{} & G\,B \end{array}$$

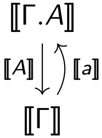
- $A \simeq B$ means isomorphism (for sets).
- This is provable in HoTT.
- It follows from univalence + J.

Summary

- The set-level fragment of HoTT can be interpreted using the groupoid model (Hofmann & Streicher).
- This interpretation also gives rise to a univalent, truncated universe of sets (but it doesn't classify hsets).
- Can we replace groupoids by categories?
- Yes, but we need to take care of polarities.
- And some places we do need groupoids, hence we need an operation calculating the groupoid associated to a category (the core).
- I am going to derive a type theory guided by the semantics.

The category with families of categories

Contexts	$\mathrm{Con} : \mathbf{Set}$	$\Gamma : \mathrm{Con}$	$[\![\Gamma]\!] : \mathbf{Cat}$
Types	$\mathrm{Ty} : \mathrm{Con} \to \mathbf{Set}$	$A : \mathrm{Ty}\,\Gamma$	$[\![A]\!] : [\![\Gamma]\!] \to \mathbf{Cat}$
Terms	$\mathrm{Tm} : (\Gamma : \mathrm{Con}) \to \mathrm{Ty}\,\Gamma \to \mathbf{Set}$	$a : \mathrm{Tm}\,\Gamma\,A$	$[\![a]\!] : \ldots$
Subst	$\mathrm{Tms} : \mathrm{Con} \to \mathrm{Con} \to \mathbf{Set}$	$\gamma : \mathrm{Tms}\,\Gamma\,\Delta$	$[\![\gamma]\!] : [\![\Gamma]\!] \to [\![\Delta]\!]$



Operations on contexts

$$\frac{}{\bullet : \mathrm{Con}} \qquad \frac{A : \mathrm{Ty}\,\Gamma}{\Gamma.A : \mathrm{Con}}$$

$$\begin{aligned} |[\![\bullet]\!]| &= \mathbf{1} \\ [\![\bullet]\!](x, y) &= \mathbf{1} \end{aligned}$$

$$\begin{aligned} |[\![\Gamma.A]\!]| &= (x : |[\![\Gamma]\!]|) \times |[\![A]\!]\,x| \\ [\![\Gamma.A]\!]((x, a), (y, b)) &= (f : [\![\Gamma]\!](x, y)) \times ([\![A]\!]\,y)([\![A]\!]\,f\,a, b) \end{aligned}$$

Grothendieck construction

Opposites

$$\frac{\Gamma : \mathrm{Con}}{\Gamma^{\mathrm{op}} : \mathrm{Con}} \qquad \frac{A : \mathrm{Ty}\,\Gamma}{A^{\mathrm{op}} : \mathrm{Ty}\,\Gamma}$$

$$[\![\Gamma^{\mathrm{op}}]\!] = [\![\Gamma]\!]^{\mathrm{op}}$$
$$[\![A^{\mathrm{op}}]\!]\,x = ([\![A]\!]\,x)^{\mathrm{op}}$$

- Note that $\_^{\mathrm{op}} : \mathbf{Cat} \to \mathbf{Cat}$ is covariant!
- But what is $(\Gamma.A)^{\mathrm{op}}$?
- It cannot be $\Gamma^{\mathrm{op}}.A^{\mathrm{op}}$

Opfibrations

$$\frac{A : \mathrm{Ty}\,\Gamma^{\mathrm{op}}}{\Gamma.^{\mathrm{op}} A : \mathrm{Con}}$$

$$|[\![\Gamma.^{\mathrm{op}} A]\!]| = (x : |[\![\Gamma|]\!]) \times |[\![A]\!]\,x|$$
$$[\![\Gamma.^{\mathrm{op}} A]\!]((x,a),(y,b)) = (f : [\![\Gamma]\!](x,y)) \times ([\![A]\!]\,x)(a, [\![A]\!]\,f\,b)$$

$$(\Gamma.A)^{\mathrm{op}} = \Gamma^{\mathrm{op}}.^{\mathrm{op}} A^{\mathrm{op}}$$

Σ-types, undirected

$$\frac{A : \mathrm{Ty}\,\Gamma \quad B : \mathrm{Ty}\,\Gamma.A}{\Sigma\,A\,B : \mathrm{Ty}\,\Gamma}$$

$$\Gamma.A.B \cong \Gamma.(\Sigma\,A\,B)$$

On objects:

$$(\Sigma\,A\,B)\,x = (A\,x).(B\,x)$$

Σ-types with polarities

$$\frac{A : \mathrm{Ty}\,\Gamma^{s} \quad B : \mathrm{Ty}\,\Gamma.^{s}A}{\Sigma^{s}\,A\,B : \mathrm{Ty}\,\Gamma}$$

$$\Gamma.^{t}A.^{s}B \cong \Gamma.^{t}(\Sigma^{s}\,A\,B)$$

On objects:

$$(\Sigma^{s}\,A\,B)\,x = (A\,x).^{s}(B\,x)$$

$$(\Sigma\,A\,B)^{\mathrm{op}} = \Sigma^{\mathrm{op}}\,A^{\mathrm{op}}\,B^{\mathrm{op}}$$

Π-types, undirected

$$\frac{A : \mathrm{Ty}\,\Gamma \quad B : \mathrm{Ty}\,\Gamma.A}{\Pi\,A\,B : \mathrm{Ty}\,\Gamma}$$

$$\mathrm{Tm}\,\Gamma.A\,B \cong \mathrm{Tm}\,\Gamma\,(\Pi\,A\,B)$$

Π-types with polarities

$$\frac{A : \mathrm{Ty}\,\Gamma^{\mathrm{op}} \quad B : \mathrm{Ty}\,\Gamma.^{\mathrm{op}}A^{s}}{\Pi^{s}\,A\,B : \mathrm{Ty}\,\Gamma}$$

$$\mathrm{Tm}\,\Gamma.^{\mathrm{op}}A^{s}\,B \cong \mathrm{Tm}\,\Gamma\,(\Pi^{s}\,A\,B)$$

$$(\Pi\,A\,B)^{\mathrm{op}} = \Pi^{\mathrm{op}}\,A^{\mathrm{op}}\,B^{\mathrm{op}}$$

The universe of sets

$$\frac{}{\mathrm{U} : \mathrm{Ty}\,\Gamma} \qquad \frac{a : \mathrm{Tm}\,\Gamma\,\mathrm{U}}{\mathrm{El}\,a : \mathrm{Ty}\,\Gamma}$$

$$|[\![\mathrm{U}]\!]\,x| = \mathbf{Set}$$
$$([\![\mathrm{U}]\!]\,x)(A, B) = A \to B$$

$$|[\![\mathrm{El}\,a]\!]\,x| = [\![a]\!]\,x$$
$$([\![\mathrm{El}\,a]\!]\,x)(y, z) = (y = z)$$

The hom type

$$\frac{A : \mathrm{Ty}\,\Gamma}{\_ \sqsubseteq_A \_ : \mathrm{Tm}\,\Gamma\,(A^{\mathrm{op}} \to A \to \mathrm{U})}$$

$$[\![a \sqsubseteq_A b]\!]\,x = A\,x(a\,x, b\,x)$$

- But what about id (aka refl)?
- We would like to say

  $$\mathrm{id} : \Pi a : A.a \sqsubseteq_A a$$

  but this doesn't type check!

The core type

$$
\frac{A : \mathrm{Ty}\,\Gamma}{\bar{A} : \mathrm{Ty}\,\Gamma}
\qquad
\frac{a : \mathrm{Tm}\,\Gamma\,A}{\bar{a} : \mathrm{Tm}\,\Gamma\,\bar{A}}
\qquad
\frac{\begin{array}{c} f : \mathrm{Tm}\,\Gamma\,(a \sqsubseteq_A b) \\ f^{\mathrm{op}} : \mathrm{Tm}\,\Gamma\,(b \sqsubseteq_A a) \\ l : \mathrm{Tm}\,\Gamma\,(f \circ f^{\mathrm{op}} \sqsubseteq \mathrm{id}_a) \\ r : \mathrm{Tm}\,\Gamma\,(f^{\mathrm{op}} \circ f \sqsubseteq \mathrm{id}_a) \end{array}}{\overline{f, f^{\mathrm{op}}, l, r} : \mathrm{Tm}\,\Gamma\,(\bar{a} \sqsubseteq_{\bar{A}} \bar{b})}
$$

$$
\frac{a : \mathrm{Tm}\,\Gamma\,\bar{A}}{\underline{a} : \mathrm{Tm}\,\Gamma\,A^s}
\qquad
\frac{f : \mathrm{Tm}\,\Gamma\,(a \sqsubseteq_{\bar{A}} b)}{\underline{f} : \mathrm{Tm}\,\Gamma\,(\underline{a} \sqsubseteq_A \underline{b})}
$$

Directed Path induction (J)

$$\mathrm{id} : \mathrm{Tm}\,\Gamma\,(\Pi a : \bar{A}.\underline{a} \sqsubseteq_A \underline{a})$$

$$\frac{\begin{array}{c} M : \mathrm{Tm}\,\Gamma\,(\Pi a : A^{\mathrm{op}}, b : A.a \sqsubseteq_A b \to \mathrm{U}) \\ m_{\mathrm{id}} : \mathrm{Tm}\,\Gamma\,(\Pi a : \bar{A}.M\,\underline{a}\,\underline{a}\,(\mathrm{id}\,a)) \end{array}}{J\,M\,m_{\mathrm{id}} : \mathrm{Tm}\,\Gamma\,(\Pi a : A^{\mathrm{op}}, b : A, f : a \sqsubseteq b.M\,a\,b\,f)}$$

The homtypes of sets

Homtypes of sets are symmetric

$$\frac{a : \mathrm{Tm}\,\Gamma\, A}{\overline{\mathrm{El}\, a} \simeq \mathrm{El}\, a}$$

Homtypes of sets are proof irrelevant

$$\frac{a : \mathrm{Tm}\,\Gamma\, A}{\mathrm{K}_a : \mathrm{Tm}\,\Gamma\,(\Pi a : \bar{A}, p : \underline{a} \sqsubseteq_A \underline{a}.p \sqsubseteq \mathrm{id}\, a)}$$

Directed univalence

$$\mathrm{coe} : \mathrm{Tm}\,\Gamma\,(\Pi a, b : \mathrm{U}.a \sqsubseteq_{\mathrm{U}} b \to (\mathrm{El}\, a \to \mathrm{El}\, b))$$

coe is an isomorphism.

Everything is natural, provably!

$$F, G : \underline{\mathbf{Set}} \to \underline{\mathbf{Set}}$$
$$\alpha : \Pi_{A:\mathbf{Set}} F\,A \to G\,A$$
$$f : A \to B$$

$$\begin{array}{ccc} F\,A & \xrightarrow{\alpha_A} & G\,A \\ {\scriptstyle F\,f}\downarrow & & \downarrow{\scriptstyle G\,f} \\ F\,B & \xrightarrow[\alpha_B]{} & G\,B \end{array}$$

- It follows from directed univalence + directed J.

Further work

- Filippo is formalizing the calculus and its semantics in Agda.
- What is the relation to logical relations?
- Can we do higher categories (full directed HoTT)?